

Manifold Learning for Signal and Visual Processing

Lecture 3: Introduction to Graphs, Graph Matrices, and Graph Embeddings

Radu Horaud

INRIA Grenoble Rhone-Alpes, France

Radu.Horaud@inrialpes.fr

<http://perception.inrialpes.fr/>

Outline of Lecture 3

- What is spectral graph theory?
- Some graph notation and definitions
- The adjacency matrix
- Laplacian matrices
- Spectral graph embedding

Material for this lecture

- F. R. K. Chung. Spectral Graph Theory. 1997. (Chapter 1)
- M. Belkin and P. Niyogi. Laplacian Eigenmaps for Dimensionality Reduction and Data Representation. Neural Computation, 15, 1373–1396 (2003).
- U. von Luxburg. A Tutorial on Spectral Clustering. Statistics and Computing, 17(4), 395–416 (2007). ([An excellent paper](#))
- Software:
<http://open-specmatch.gforge.inria.fr/index.php>.
Computes, among others, Laplacian embeddings of very large graphs.

Spectral graph theory at a glance

- The *spectral graph theory* studies the properties of graphs via the eigenvalues and eigenvectors of their associated graph matrices: the *adjacency matrix*, the *graph Laplacian* and their variants.
- These matrices have been extremely well studied from an algebraic point of view.
- The Laplacian allows a natural link between discrete representations (graphs), and continuous representations, such as metric spaces and manifolds.
- Laplacian embedding consists in representing the vertices of a graph in the space spanned by the smallest eigenvectors of the Laplacian – *A geodesic distance on the graph becomes a spectral distance in the embedded (metric) space.*

Spectral graph theory and manifold learning

- First we construct a graph from $\mathbf{x}_1, \dots, \mathbf{x}_n \in \mathbb{R}^D$
- Then we compute the d **smallest** eigenvalue-eigenvector pairs of the graph Laplacian
- Finally we represent the data in the \mathbb{R}^d space spanned by the corresponding orthonormal eigenvector basis. The choice of the dimension d of the embedded space is not trivial.
- Paradoxically, d may be **larger** than D in many cases!

Basic graph notations and definitions

We consider *simple graphs* (no multiple edges or loops),
 $\mathcal{G} = \{\mathcal{V}, \mathcal{E}\}$:

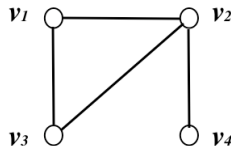
- $\mathcal{V}(\mathcal{G}) = \{v_1, \dots, v_n\}$ is called the *vertex set* with $n = |\mathcal{V}|$;
- $\mathcal{E}(\mathcal{G}) = \{e_{ij}\}$ is called the *edge set* with $m = |\mathcal{E}|$;
- An edge e_{ij} connects vertices v_i and v_j if they are adjacent or neighbors. One possible notation for adjacency is $v_i \sim v_j$;
- The number of neighbors of a node v is called the *degree* of v and is denoted by $d(v)$, $d(v_i) = \sum_{v_i \sim v_j} e_{ij}$. If all the nodes of a graph have the same degree, the graph is *regular*; The nodes of an *Eulerian* graph have even degree.
- A graph is *complete* if there is an edge between every pair of vertices.

The adjacency matrix of a graph

- For a graph with n vertices, the entries of the $n \times n$ adjacency matrix are defined by:

$$\mathbf{A} := \begin{cases} A_{ij} = 1 & \text{if there is an edge } e_{ij} \\ A_{ij} = 0 & \text{if there is no edge} \\ A_{ii} = 0 \end{cases}$$

$$\mathbf{A} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$



Eigenvalues and eigenvectors

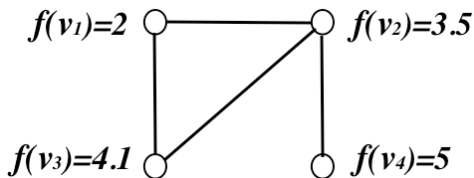
- \mathbf{A} is a real-symmetric matrix: it has n real eigenvalues and its n real eigenvectors form an orthonormal basis.
- Let $\{\lambda_1, \dots, \lambda_i, \dots, \lambda_r\}$ be the set of *distinct* eigenvalues.
- The eigenspace S_i contains the eigenvectors associated with λ_i :

$$S_i = \{\mathbf{x} \in \mathbb{R}^n \mid \mathbf{A}\mathbf{x} = \lambda_i\mathbf{x}\}$$

- For real-symmetric matrices, the algebraic multiplicity is equal to the geometric multiplicity, for all the eigenvalues.
- The dimension of S_i (geometric multiplicity) is equal to the multiplicity of λ_i .
- If $\lambda_i \neq \lambda_j$ then S_i and S_j are mutually orthogonal.

Real-valued functions on graphs

- We consider real-valued functions on the set of the graph's vertices, $f : \mathcal{V} \longrightarrow \mathbb{R}$. Such a function assigns a real number to each graph node.
- f is a vector indexed by the graph's vertices, hence $f \in \mathbb{R}^n$.
- **Notation:** $f = (f(v_1), \dots, f(v_n)) = (f_1, \dots, f_n)$.
- The eigenvectors of the adjacency matrix, $\mathbf{A}x = \lambda x$, can be viewed as *eigenfunctions*.



Matrix \mathbf{A} as an operator and quadratic form

- The adjacency matrix can be viewed as an operator

$$\mathbf{g} = \mathbf{A}\mathbf{f}; g(i) = \sum_{i \sim j} f(j)$$

- It can also be viewed as a quadratic form:

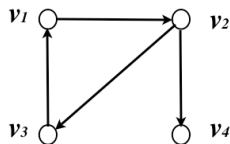
$$\mathbf{f}^\top \mathbf{A} \mathbf{f} = \sum_{e_{ij}} f(i)f(j)$$

The incidence matrix of a graph

- Let each edge in the graph have an arbitrary but fixed orientation;
- The incidence matrix of a graph is a $|\mathcal{E}| \times |\mathcal{V}|$ ($m \times n$) matrix defined as follows:

$$\nabla := \begin{cases} \nabla_{ev} = -1 & \text{if } v \text{ is the initial vertex of edge } e \\ \nabla_{ev} = 1 & \text{if } v \text{ is the terminal vertex of edge } e \\ \nabla_{ev} = 0 & \text{if } v \text{ is not in } e \end{cases}$$

$$\nabla = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & 0 & -1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & +1 \end{bmatrix}$$



The incidence matrix: A discrete differential operator

- The mapping $f \longrightarrow \nabla f$ is known as the *co-boundary mapping* of the graph.
- $(\nabla f)(e_{ij}) = f(v_j) - f(v_i)$

$$\begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & 0 & -1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & +1 \end{bmatrix} \begin{pmatrix} f(1) \\ f(2) \\ f(3) \\ f(4) \end{pmatrix} = \begin{pmatrix} f(2) - f(1) \\ f(1) - f(3) \\ f(3) - f(2) \\ f(4) - f(2) \end{pmatrix}$$

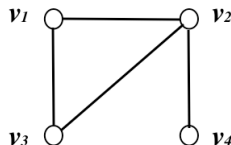
The Laplacian matrix of a graph

- $\mathbf{L} = \nabla^\top \nabla$
- $(\mathbf{L}\mathbf{f})(v_i) = \sum_{v_j \sim v_i} (f(v_i) - f(v_j))$
- Connection between the Laplacian and the adjacency matrices:

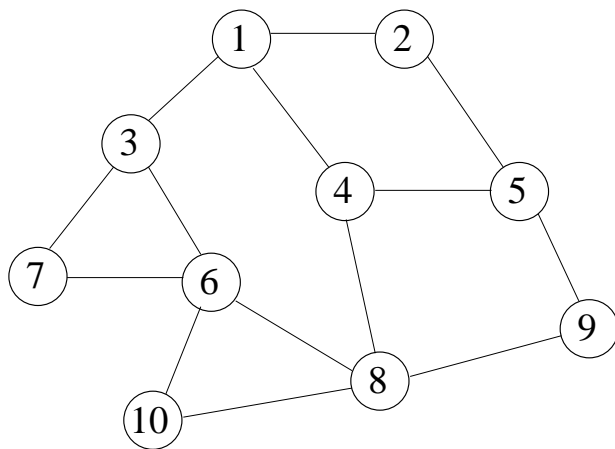
$$\mathbf{L} = \mathbf{D} - \mathbf{A}$$

- The degree matrix: $\mathbf{D} := D_{ii} = d(v_i)$.

$$\mathbf{L} = \begin{bmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 2 & 0 \\ 0 & -1 & 0 & 1 \end{bmatrix}$$



Example: A graph with 10 nodes



The adjacency matrix

$$\mathbf{A} = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \end{bmatrix}$$

The Laplacian matrix

$$\mathbf{L} = \begin{bmatrix} 3 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 2 & 0 & 0 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 3 & 0 & 0 & -1 & -1 & 0 & 0 & 0 \\ -1 & 0 & 0 & 3 & -1 & 0 & 0 & -1 & 0 & 0 \\ 0 & -1 & 0 & -1 & 3 & 0 & 0 & 0 & -1 & 0 \\ 0 & 0 & -1 & 0 & 0 & 4 & -1 & -1 & 0 & -1 \\ 0 & 0 & -1 & 0 & 0 & -1 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 & 0 & -1 & 0 & 4 & -1 & -1 \\ 0 & 0 & 0 & 0 & -1 & 0 & 0 & -1 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & -1 & 0 & -1 & 0 & 2 \end{bmatrix}$$

The Eigenvalues of this Laplacian

$$\mathbf{\Lambda} = \begin{bmatrix} 0.0000 & 0.7006 & 1.1306 & 1.8151 & 2.4011 \\ 3.0000 & 3.8327 & 4.1722 & 5.2014 & 5.7462 \end{bmatrix}$$

Matrices of an Undirected Weighted Graph

- We consider *undirected weighted graphs*; Each edge e_{ij} is weighted by $w_{ij} > 0$. We obtain:

$$\mathbf{\Omega} := \begin{cases} \Omega_{ij} = w_{ij} & \text{if there is an edge } e_{ij} \\ \Omega_{ij} = 0 & \text{if there is no edge} \\ \Omega_{ii} = 0 \end{cases}$$

- The degree matrix: $\mathbf{D} = \sum_{i \sim j} w_{ij}$

The Laplacian on an undirected weighted graph

- $\mathbf{L} = \mathbf{D} - \mathbf{\Omega}$
- The Laplacian as an operator:

$$(\mathbf{L}\mathbf{f})(v_i) = \sum_{v_j \sim v_i} w_{ij}(f(v_i) - f(v_j))$$

- As a quadratic form:

$$\mathbf{f}^\top \mathbf{L} \mathbf{f} = \frac{1}{2} \sum_{e_{ij}} w_{ij} (f(v_i) - f(v_j))^2$$

- \mathbf{L} is symmetric and positive semi-definite $\leftrightarrow w_{ij} \geq 0$.
- \mathbf{L} has n non-negative, real-valued eigenvalues:
 $0 = \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$.

Other adjacency matrices

- The *normalized weighted adjacency matrix*

$$\Omega_N = \mathbf{D}^{-1/2} \Omega \mathbf{D}^{-1/2}$$

- The *transition* matrix of the Markov process associated with the graph:

$$\Omega_R = \mathbf{D}^{-1} \Omega = \mathbf{D}^{-1/2} \Omega_N \mathbf{D}^{1/2}$$

Several Laplacian matrices

- The *unnormalized Laplacian* which is also referred to as the *combinatorial Laplacian* \mathbf{L}_C ,
- the *normalized Laplacian* \mathbf{L}_N , and
- the *random-walk Laplacian* \mathbf{L}_R also referred to as the *discrete Laplace operator*.

We have:

$$\mathbf{L}_C = \mathbf{D} - \mathbf{\Omega}$$

$$\mathbf{L}_N = \mathbf{D}^{-1/2} \mathbf{L}_C \mathbf{D}^{-1/2} = \mathbf{I} - \mathbf{\Omega}_N$$

$$\mathbf{L}_R = \mathbf{D}^{-1} \mathbf{L}_C = \mathbf{I} - \mathbf{\Omega}_R$$

Relationships between all these matrices

$$\mathbf{L}_C = \mathbf{D}^{1/2} \mathbf{L}_N \mathbf{D}^{1/2} = \mathbf{D} \mathbf{L}_R$$

$$\mathbf{L}_N = \mathbf{D}^{-1/2} \mathbf{L}_C \mathbf{D}^{-1/2} = \mathbf{D}^{1/2} \mathbf{L}_R \mathbf{D}^{-1/2}$$

$$\mathbf{L}_R = \mathbf{D}^{-1/2} \mathbf{L}_N \mathbf{D}^{1/2} = \mathbf{D}^{-1} \mathbf{L}_C$$

Some spectral properties of the Laplacians

Laplacian	Null space	Eigenvalues	Eigenvectors
$\mathbf{L}_C = \mathbf{U}\mathbf{\Lambda}\mathbf{U}^\top$	$\mathbf{u}_1 = \mathbf{1}$	$0 = \lambda_1 < \lambda_2 \leq \dots \leq \lambda_n \leq 2 \max_i(d_i)$	$\mathbf{u}_{i>1}^\top \mathbf{1} = 0,$ $\mathbf{u}_i^\top \mathbf{u}_j = \delta_{ij}$
$\mathbf{L}_N = \mathbf{W}\mathbf{\Gamma}\mathbf{W}^\top$	$\mathbf{w}_1 = \mathbf{D}^{1/2}\mathbf{1}$	$0 = \gamma_1 < \gamma_2 \leq \dots \leq \gamma_n \leq 2$	$\mathbf{w}_{i>1}^\top \mathbf{D}^{1/2}\mathbf{1} = 0,$ $\mathbf{w}_i^\top \mathbf{w}_j = \delta_{ij}$
$\mathbf{L}_R = \mathbf{T}\mathbf{T}\mathbf{T}^{-1}$ $\mathbf{T} = \mathbf{D}^{-1/2}\mathbf{W}$	$\mathbf{t}_1 = \mathbf{1}$	$0 = \gamma_1 < \gamma_2 \leq \dots \leq \gamma_n \leq 2$	$\mathbf{t}_{i>1}^\top \mathbf{D}\mathbf{1} = 0,$ $\mathbf{t}_i^\top \mathbf{D}\mathbf{t}_j = \delta_{ij}$

Spectral properties of adjacency matrices

From the relationship between the normalized Laplacian and adjacency matrix: $\mathbf{L}_N = \mathbf{I} - \mathbf{\Omega}_N$ one can see that their eigenvalues satisfy:

$$\gamma = 1 - \psi$$

Adjacency matrix	Eigenvalues	Eigenvectors
$\mathbf{\Omega}_N = \mathbf{W}\mathbf{\Psi}\mathbf{W}^\top$, $\mathbf{\Psi} = \mathbf{I} - \mathbf{\Gamma}$	$-1 \leq \psi_n \leq \dots \leq \psi_2 < \psi_1 = 1$	$\mathbf{w}_i^\top \mathbf{w}_j = \delta_{ij}$
$\mathbf{\Omega}_R = \mathbf{T}\mathbf{\Psi}\mathbf{T}^{-1}$	$-1 \leq \psi_n \leq \dots \leq \psi_2 < \psi_1 = 1$	$\mathbf{t}_i^\top \mathbf{D}\mathbf{t}_j = \delta_{ij}$

Eigenvalue and Eigenvectors of the Normalized and Random Laplacians

- Eigenvalues of the normalized adjacent matrix:

$$1 = \psi_1 \geq \psi_2 \geq \dots \geq \psi_n \geq -1$$

- The largest eigenvalue-eigenvector pair:
($\psi_1 = 1, \mathbf{w}_1 = \mathbf{D}^{1/2}\mathbf{1}$)
- The estimation of the smallest non null eigenvalue-eigenvector pairs of \mathbf{L}_N involves the shifted inverse power method.
- The second, third, etc., largest eigenvalue-eigenvector pair of $\mathbf{\Omega}_N$ can be obtained with the direct power method and deflation:

$$\tilde{\mathbf{\Omega}}_N = \mathbf{\Omega}_N - \mathbf{w}_1 \mathbf{w}_1^\top$$

.

- **Remark:** Sparsity is lost by deflation!

The Laplacian of a graph with one connected component

- $\mathbf{L}\mathbf{u} = \lambda\mathbf{u}$.
- $\mathbf{L}\mathbf{1} = \mathbf{0}$, $\lambda_1 = 0$ is the smallest eigenvalue.
- The *one* vector: $\mathbf{1} = (1 \dots 1)^\top$.
- $0 = \mathbf{u}^\top \mathbf{L}\mathbf{u} = \sum_{i,j=1}^n w_{ij}(u(i) - u(j))^2$.
- If any two vertices are connected by a path, then $\mathbf{u} = (u(1), \dots, u(n))$ needs to be constant at all vertices such that the quadratic form vanishes. Therefore, a graph with one connected component has the constant vector $\mathbf{u}_1 = \mathbf{1}$ as the only eigenvector with eigenvalue 0.

A graph with $k > 1$ connected components

- Each connected component has an associated Laplacian.
Therefore, we can write matrix \mathbf{L} as a *block diagonal matrix*:

$$\mathbf{L} = \begin{bmatrix} \mathbf{L}_1 & & \\ & \ddots & \\ & & \mathbf{L}_k \end{bmatrix}$$

- The spectrum of \mathbf{L} is given by the union of the spectra of \mathbf{L}_i .
- Each block corresponds to a connected component, hence each matrix \mathbf{L}_i has an eigenvalue 0 with multiplicity 1.
- The spectrum of \mathbf{L} is given by the union of the spectra of \mathbf{L}_i .
- The eigenvalue $\lambda_1 = 0$ has multiplicity k .

The eigenspace of $\lambda_1 = 0$ with multiplicity k

- The eigenspace corresponding to $\lambda_1 = \dots = \lambda_k = 0$ is spanned by the k mutually orthogonal vectors:

$$\mathbf{u}_1 = \mathbf{1}_{L_1}$$

$$\dots$$

$$\mathbf{u}_k = \mathbf{1}_{L_k}$$

- with $\mathbf{1}_{L_i} = (0000111110000)^\top \in \mathbb{R}^n$
- These vectors are the *indicator vectors* of the graph's connected components.
- Notice that $\mathbf{1}_{L_1} + \dots + \mathbf{1}_{L_k} = \mathbf{1}$

The Fiedler vector of the graph Laplacian

- The first non-null eigenvalue λ_{k+1} is called the Fiedler value.
- The corresponding eigenvector \mathbf{u}_{k+1} is called the Fiedler vector.
- The multiplicity of the Fiedler eigenvalue depends on the graph's structure and it is difficult to analyse.
- The Fiedler value is the *algebraic connectivity of a graph*, the further from 0, the more connected.
- The Fiedler vector has been extensively used for *spectral bi-partitioning*
- Theoretical results are summarized in Spielman & Teng 2007: <http://cs-www.cs.yale.edu/homes/spielman/>

Eigenvectors of the Laplacian of connected graphs

- $\mathbf{u}_1 = \mathbf{1}, \mathbf{L}\mathbf{1} = \mathbf{0}$.
- \mathbf{u}_2 is the *Fiedler vector* with multiplicity 1.
- The eigenvectors form an orthonormal basis: $\mathbf{u}_i^\top \mathbf{u}_j = \delta_{ij}$.
- For any eigenvector $\mathbf{u}_i = (\mathbf{u}_i(v_1) \dots \mathbf{u}_i(v_n))^\top$, $2 \leq i \leq n$:

$$\mathbf{u}_i^\top \mathbf{1} = 0$$

- Hence the components of \mathbf{u}_i , $2 \leq i \leq n$ satisfy:

$$\sum_{j=1}^n \mathbf{u}_i(v_j) = 0$$

- Each component is bounded by:

$$-1 < \mathbf{u}_i(v_j) < 1$$

Laplacian embedding: Mapping a graph on a line

- Map a weighted graph onto a line such that connected nodes stay as close as possible, i.e., minimize

$\sum_{i,j=1}^n w_{ij}(f(v_i) - f(v_j))^2$, or:

$$\arg \min_{\mathbf{f}} \mathbf{f}^\top \mathbf{L} \mathbf{f} \text{ with: } \mathbf{f}^\top \mathbf{f} = 1 \text{ and } \mathbf{f}^\top \mathbf{1} = 0$$

- The solution is the eigenvector associated with the smallest nonzero eigenvalue of the eigenvalue problem: $\mathbf{L} \mathbf{f} = \lambda \mathbf{f}$, namely the Fiedler vector \mathbf{u}_2 .
- Practical computation of the eigenpair λ_2, \mathbf{u}_2): the shifted inverse power method (see lecture 2).

The shifted inverse power method (from Lecture 2)

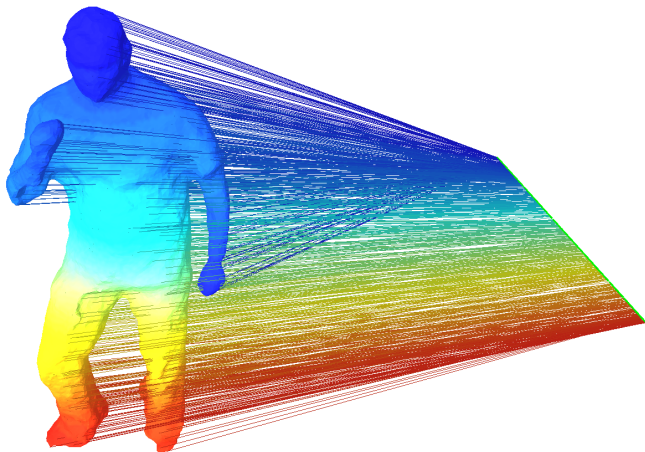
- Let's consider the matrix $\mathbf{B} = \mathbf{A} - \alpha\mathbf{I}$ as well as an eigenpair $\mathbf{A}\mathbf{u} = \lambda\mathbf{u}$.
- $(\lambda - \alpha, \mathbf{u})$ becomes an eigenpair of \mathbf{B} , indeed:

$$\mathbf{B}\mathbf{u} = (\mathbf{A} - \alpha\mathbf{I})\mathbf{u} = (\lambda - \alpha)\mathbf{u}$$

and hence \mathbf{B} is a **real symmetric** matrix with eigenpairs $(\lambda_1 - \alpha, \mathbf{u}_1), \dots, (\lambda_i - \alpha, \mathbf{u}_i), \dots, (\lambda_D - \alpha, \mathbf{u}_D)$

- If $\alpha > 0$ is chosen such that $|\lambda_j - \alpha| \ll |\lambda_i - \alpha| \forall i \neq j$ then $\lambda_j - \alpha$ becomes the smallest (in magnitude) eigenvalue.
- The inverse power method (in conjunction with the *LU decomposition* of \mathbf{B}) can be used to estimate the eigenpair $(\lambda_j - \alpha, \mathbf{u}_j)$.

Example of mapping a graph on the Fiedler vector



Laplacian embedding

- Embed the graph in a k -dimensional Euclidean space. The embedding is given by the $n \times k$ matrix $\mathbf{F} = [\mathbf{f}_1 \mathbf{f}_2 \dots \mathbf{f}_k]$ where the i -th row of this matrix – $\mathbf{f}^{(i)}$ – corresponds to the Euclidean coordinates of the i -th graph node v_i .
- We need to minimize (Belkin & Niyogi '03):

$$\arg \min_{\mathbf{f}_1 \dots \mathbf{f}_k} \sum_{i,j=1}^n w_{ij} \|\mathbf{f}^{(i)} - \mathbf{f}^{(j)}\|^2 \text{ with: } \mathbf{F}^\top \mathbf{F} = \mathbf{I}.$$

- The solution is provided by the matrix of eigenvectors corresponding to the k lowest nonzero eigenvalues of the eigenvalue problem $\mathbf{L}\mathbf{f} = \lambda\mathbf{f}$.

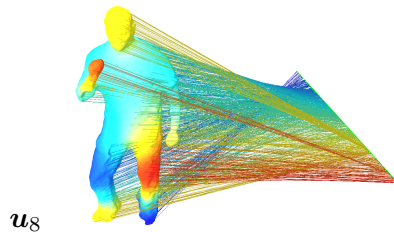
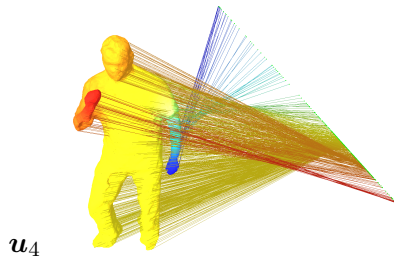
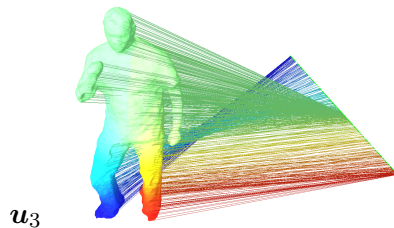
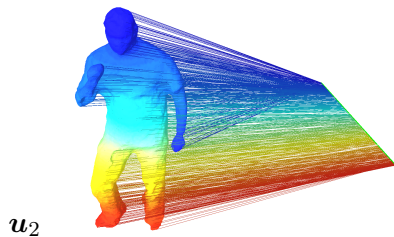
Spectral embedding using the *unnormalized* Laplacian

- Compute the eigendecomposition $\mathbf{L} = \mathbf{D} - \mathbf{\Omega}$.
- Select the k smallest non-null eigenvalues $\lambda_2 \leq \dots \leq \lambda_{k+1}$
- $\lambda_{k+2} - \lambda_{k+1} = \mathbf{eigengap}$.
- We obtain the $n \times k$ matrix $\mathbf{U} = [\mathbf{u}_2 \dots \mathbf{u}_{k+1}]$:

$$\mathbf{U} = \begin{bmatrix} \mathbf{u}_2(v_1) & \dots & \mathbf{u}_{k+1}(v_1) \\ \vdots & & \vdots \\ \mathbf{u}_2(v_n) & \dots & \mathbf{u}_{k+1}(v_n) \end{bmatrix}$$

- $\mathbf{u}_i^\top \mathbf{u}_j = \delta_{ij}$ (orthonormal vectors), hence $\mathbf{U}^\top \mathbf{U} = \mathbf{I}_k$.
- Column i ($2 \leq i \leq k+1$) of this matrix is a mapping on the eigenvector \mathbf{u}_i .

Examples of one-dimensional mappings



Euclidean L-embedding of the graph's vertices

- (Euclidean) **L**-embedding of a graph:

$$\mathbf{X} = \mathbf{\Lambda}_k^{-\frac{1}{2}} \mathbf{U}^\top = [\mathbf{x}_1 \ \dots \ \mathbf{x}_j \ \dots \ \mathbf{x}_n]$$

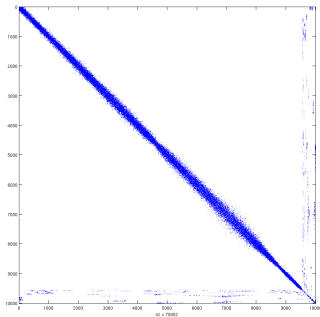
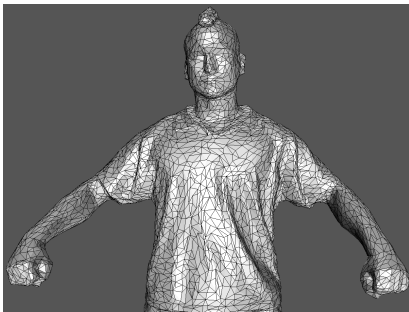
The coordinates of a vertex v_j are:

$$\mathbf{x}_j = \begin{pmatrix} \frac{\mathbf{u}_2(v_j)}{\sqrt{\lambda_2}} \\ \vdots \\ \frac{\mathbf{u}_{k+1}(v_j)}{\sqrt{\lambda_{k+1}}} \end{pmatrix}$$

- A formal justification of using this will be provided later.

The Laplacian of a mesh

A mesh may be viewed as a graph: $n = 10,000$ vertices, $m = 35,000$ edges. ARPACK finds the smallest 100 eigenpairs in 46 seconds.



Example: Shape embedding

